AMENDMENTS TO THE CLAIMS

This listing of claims replaces all prior versions, and listings, of claims in the application:

1	1. (Currently Amended) A method for use with performed by software embodied in
2	a computer-readable storage medium and executed by a computer in a database system that
3	stores a join view associated with plural base relations, the method comprising:
4	receiving modification operations that modify at least two of the base relations of
5	the join view, wherein the at least two base relations comprise a first base relation and a second
6	base relation; [[and]]
7	performing partitioning of the received modification operations by submitting at
8	least some of the modification operations operating on the first base relation to a first session,
9	and submitting at least another of the modification operations that operate on the second base
10	relation to a second session;
11	grouping the at least some of the modification operations in the first session
12	operating on the first base relation into a first transaction,
13	wherein the at least another modification operation in the second session is part of
14	a second transaction; and
15	re-ordering the received modification operations schedule the transactions to
16	avoid execution of modification operations of more than one of the at least two base relations at
17	one time in the database system.
1	2. (Currently Amended) The method of claim 1, wherein receiving the modification
2	operations comprises receiving a first modification operation to modify a first base relation of the
3	join view, and a second modification operation to modify a second base relation of the join view,
4	wherein re ordering the modification operations scheduling the transactions
5	comprises:
6	determining that the first modification operation transaction conflicts with
7	the second modification operation transaction based on the first and second modification
8	operations transactions modifying more than one base relation of the join view; and
9	selecting one of the first and second modification operations transactions
10	for execution in the database system.

I	3.	(Currently Amended) The method of claim $[[1]] \underline{z}$, wherein selecting one of the
2	first and seco	nd modification operations transactions comprises selecting the first modification
3	operation trai	nsaction,
4		the method further comprising storing the second modification operation
5	transaction in	a queue.
1	4.	(Currently Amended) The method of claim 3, further comprising waiting for the
2	first modifica	tion operation transaction to complete execution before scheduling the second
3	modification	operation transaction for operation execution

1	5. (Currently Amended) The method of claim 4, further comprising A method performed
2	by software embodied in a computer-readable storage medium and executed by a computer in a
3	database system that stores a join view associated with plural base relations, comprising:
4	receiving a first modification operation to modify a first base relation of the join
5	view, and a second modification operation to modify a second base relation of the join view; and
6	re-ordering the received modification operations to avoid execution of
7	modification operations of more than one of the first and second base relations at one time in the
8	database system,
9	wherein re-ordering the modification operations comprises:
10	determining that the first modification operation conflicts with the second
11	modification operation based on the first and second modification operations modifying more
12	than one base relation of the join view;
13	selecting the first modification operation for execution in the database
14	system;
15	storing the second modification operation in a queue;
16	waiting for the first modification operation to complete execution before
17	scheduling the second modification operation for operation;
18	receiving a third modification operation to modify the first base relation of the
19	join view , the method further comprising:
20	storing the third modification operation in the queue; and
21	scheduling the third modification operation for execution in the database system
22	ahead of the second modification operation.
1	6 _ 7 (Cancelled)

1	8.	(Currently Amended) The method of claim 1, further comprising: A method
2	performed by	software embodied in a computer-readable storage medium and executed by a
3	computer in a	database system that stores a join view associated with plural base relations,
4	comprising:	
5		receiving modification operations that modify at least two of the base relations of
6	the join view;	
7		re-ordering the received modification operations to avoid execution of
8	modification	operations of more than one of the at least two base relations at one time in the
9	database syste	em; and
10		in response to a <u>particular one of the</u> modification operation <u>operations</u> to modify
11	one of the bas	se relations, placing an exclusive lock on the one base relation, and placing a
12	predefined lo	ck on the join view,
13		the predefined lock conflicting with either each of a shared lock [[or]] and an
14	exclusive loc	k placed on the join view, but the predefined lock not conflicting with another
15	predefined lo	ck placed on the join view.
1	9.	(Currently Amended) The method of claim 1, further comprising:
2		storing pending modification operations transactions in plural queues
3	corresponding	g to respective plural sessions of the database system; and
4		selecting one of the pending modification operations transactions from the queues
5	to schedule fo	or execution in the database system based on whether the one pending modification
6	operation tran	nsaction conflicts with one or more executing modification operations transactions
7	in the databas	se system.
1	10.	(Currently Amended) The method of claim 9, further comprising determining
2	that the one p	ending modification operation transaction conflicts with the one or more executing
3	modification	operations transactions in response to determining that the one pending
4	modification-	operation transaction modifies a different one of the base relations of the join view
5	than a base re	elation of the join view modified by an executing modification operation
6	transaction.	

1	11. (Currently Amended) The method of claim 9, further comprising applying a
2	technique to prevent starvation of <u>a particular</u> one of the pending modification operations
3	transactions in response to determining that the particular one pending modification operation
4	transaction has been in one of the queues for longer than a predetermined time period.
1	12. (Currently Amended) An article comprising at least one computer-readable
2	storage medium containing instructions that when executed cause a system computer to:
3	receive modification operations that modify at least two of the base relations of a
4	join view, wherein the at least two base relations comprise a first base relation and a second base
5	relation;
6	perform partitioning of the received modification operations by submitting at least
7	some of the modification operations operating on the first base relation to a first session, and
8	submitting at least another of the modification operations that operate on a second base relation
9	to a second session;
10	group the at least some of the modification operations in the first session
11	operating on the first base relation into a first transaction,
12	wherein the at least another modification operation in the second session is part of
13	a second transaction; and
14	re order the received modification operations schedule the transactions to avoid
15	concurrent execution of modification operations transactions of more than one of the at least two
16	base relations of the join view.

1	13. (Currently Amended) The article of claim 12, wherein receiving the modification
2	operations comprises receiving a first modification operation to modify a first base relation of the
3	join view, and a second modification operation to modify a second base relation of the join view,
4	wherein re-ordering the modification operations scheduling the transactions
5	comprises:
6	determining that the first modification operation transaction conflicts with
7	the second modification operation transaction based on the first and second modification
8	operations transactions modifying more than one base relation of the join view; and
9	selecting one of the first and second modification operations transactions
10	for execution in the database system.
1	14. (Currently Amended) The article of claim 13, wherein selecting one of the first
2	and second modification operations transactions comprises selecting the first modification
3	operation transaction,
4	the instructions when executed causing the system to further store the second
5	modification operation transaction in a queue.
1	15. (Currently Amended) The article of claim 14, wherein the instructions when
2	executed cause the system computer to wait for the first modification operation transaction to
3	complete execution before scheduling the second modification operation transaction for
4	operation execution.
1	16. – 17. (Cancelled)

1	18.	(Currently Amended) The article of claim 12, wherein the instructions when
2	executed cau	se the system computer to:
3		in response to a <u>particular one of the</u> modification operation <u>operations</u> to modify
4	one of the ba	se relations, place an exclusive lock on the one base relation, and place a predefined
5	lock on the jo	oin view,
6		the predefined lock conflicting with either each of a shared lock [[or]] and an
7	exclusive loc	k placed on the join view, but the predefined lock not conflicting with another
8	predefined lo	ock placed on the join view.
1	19.	(Cancelled)
1	20.	(Currently Amended) The <u>first</u> system of claim [[19]] <u>22</u> , wherein the controller
2	is adapted to	identify the modification operations on the second base relation as conflicting with
3	the modificat	ion operations on the first base relation in response to determining that the
4	modification	operations on the second base relation are modifying a different base relation of the
5	join view tha	n the modification operations on the first base relation.
1	21.	(Currently Amended) The first system of claim [[20]] 22, wherein the system
2	comprises a	First system, and wherein the controller is adapted to open plural sessions with [[a]]
3	the database	system that is separate from the first system,
4		the controller to further:
5		determine that the identify modification operations on the first base
6	relation that	modify distinct portions of the first base relation; and
7		in response to determining that the modification operations on the first
8	base-relation	modify distinct portions of the first base relation, submit the identified modification
9	operations [[on]] that modify distinct portions of the first base relation through different sessions
10	for concurren	nt execution in the database system.

1	22. (Currently Amended) The system of claim 20, wherein the system comprises a
2	first system, A first system comprising:
3	a controller having one or more processors to:
4	receive modification operations to modify plural base relations of a join
5	view, the modification operations comprising modification operations to modify a first base
6	relation of the join view, and modification operations to modify a second base relation of the join
7	view; and
8	re-order the received modification operations to avoid concurrent
9	execution of modification operations of more than one of the plural base relations of the join
10	view,
11	the re-ordering to cause modification operations on the first base relation
12	of the join view to be scheduled for execution, and to cause modification operations on the
13	second base relation to be queued for execution after completion of the modification operations
14	on the first base relation,
15	wherein certain of the modification operations on the first base relation
16	eomprises comprise modification operations of a first tuple set of one or more tuples of the first
17	base relation, and wherein the controller is adapted to:
18	group the modification operations on the first tuple set of one or more
19	tuples of the first base relation into a transaction; and
20	submit the transaction to a database system separate from the first system
21	for execution.
1	23. (Currently Amended) The <u>first</u> system of claim [[19]] <u>22</u> , wherein the controller
2	comprises a load utility to submit the modification operations to [[a]] the database system.
1	24. (Currently Amended) The <u>first</u> system of claim 23, wherein the load utility
2	comprises a continuous load utility.

- 1 25. (Currently Amended) The <u>first</u> system of claim 23, wherein the load utility comprises a first load utility, and the controller comprises a second load utility to concurrently
- 3 submit other modification operations to the database system.
- 1 26. (Currently Amended) The <u>first</u> system of claim 25, further comprising plural
- 2 platforms on which corresponding first and second load utilities are executable.